

FROM PROGRAMMABLE LOGIC CONTROL (PLC) TO DISCRETE EVENT SYSTEMS (DES)

C. G. Cassandras

Dept. of Manufacturing Engineering

Boston University

Boston, MA 02215

cgc@bu.edu

<http://vita.bu.edu/cgc>



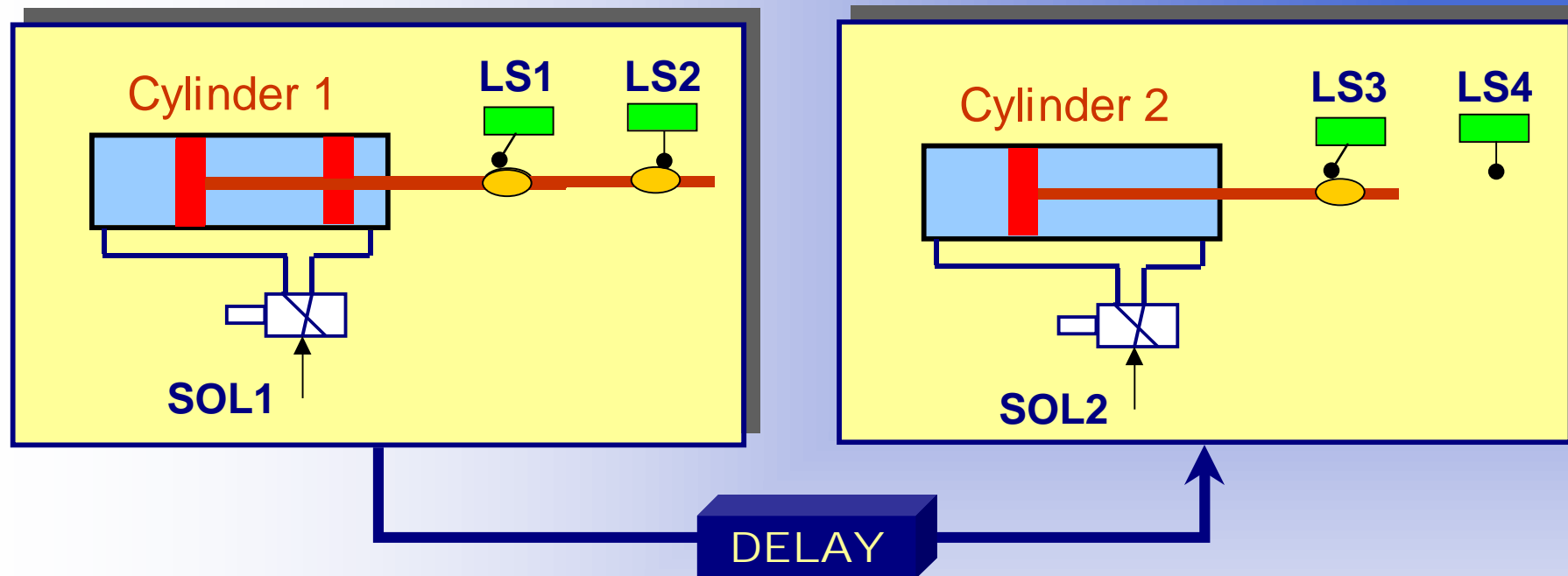
OUTLINE

- A MOTIVATING EXAMPLE
- DES MODELING:
STATE AUTOMATA
(STATE MACHINES)
- SUPERVISORY CONTROL
- ~~RA~~ ~~PP~~ ~~LE~~ ~~ARE~~ ~~IN~~ ~~ON~~ ~~FIGURATION~~ ~~IO~~

TECHNIQUES



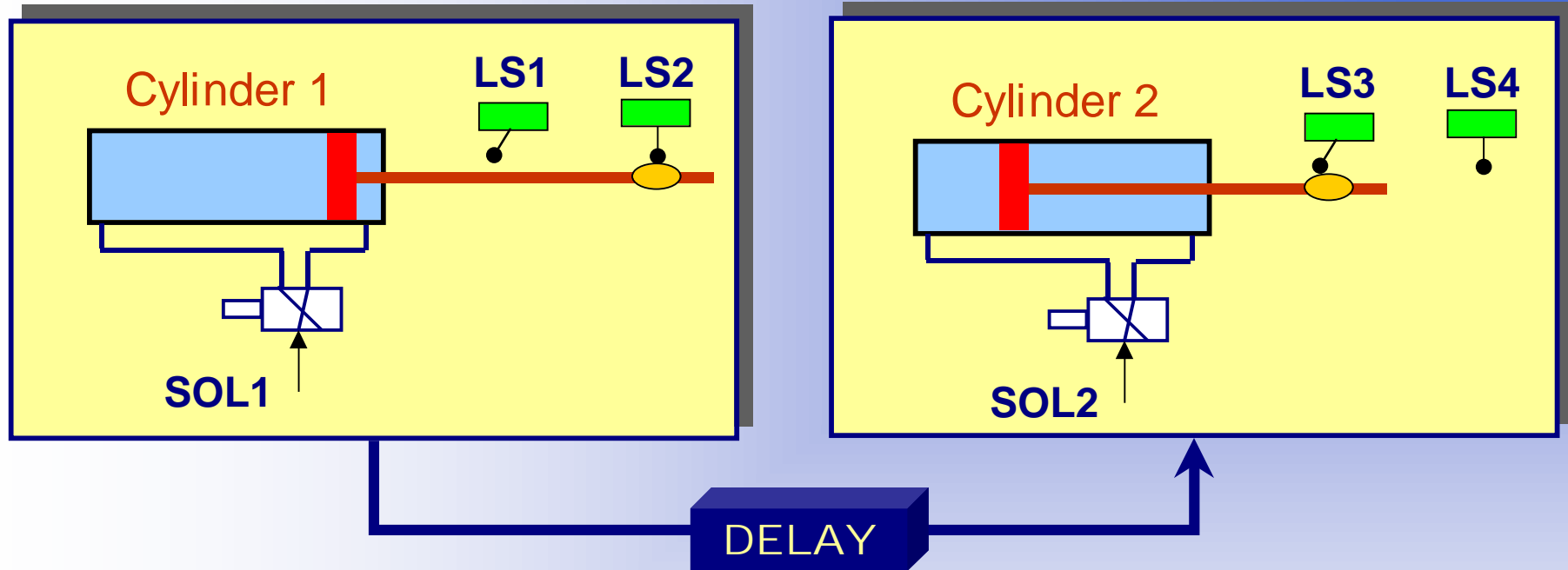
A MOTIVATING EXAMPLE



- **START** initiates cycle with **LS1** and **LS3** ON
- **SOL1** activated, piston 1 moves right
- **LS2** triggered, **SOL1** deactivated, piston 1 moves left
- Time delay of 1 sec.
- **SOL2** activated, piston 2 moves right
- **LS4** triggered, **SOL2** deactivated, piston 2 moves left
- **STOP** resets everything to rest



EXAMPLE – PLC APPROACH



Cylinder 1 control: $SOL1 = (SOL1 + START \cdot LS1 \cdot LS3) \cdot \overline{LS2} \cdot \overline{STOP}$

ACTIVATE

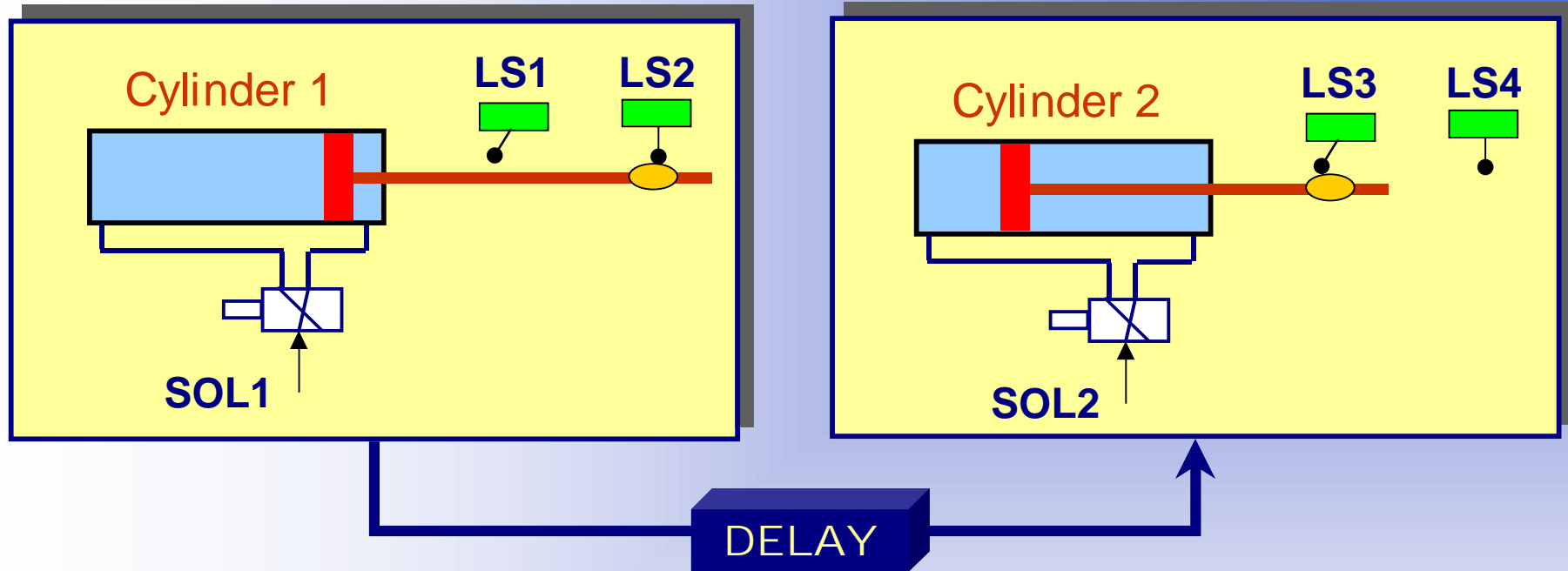
DEACTIVAT

NOTE: Upon deactivation of **SOL1**, the system looks **IDENTICAL** as at rest, yet it must know to activate the **DELAY** function.



EXAMPLE - PLC APPROACH

CONTINUED



Solution: Introduce some memory...

$$DONE = (DONE + LS2) \cdot \overline{LS4} \cdot \overline{STOP}$$

INDICATES CYLINDER1
CYCLE DONE

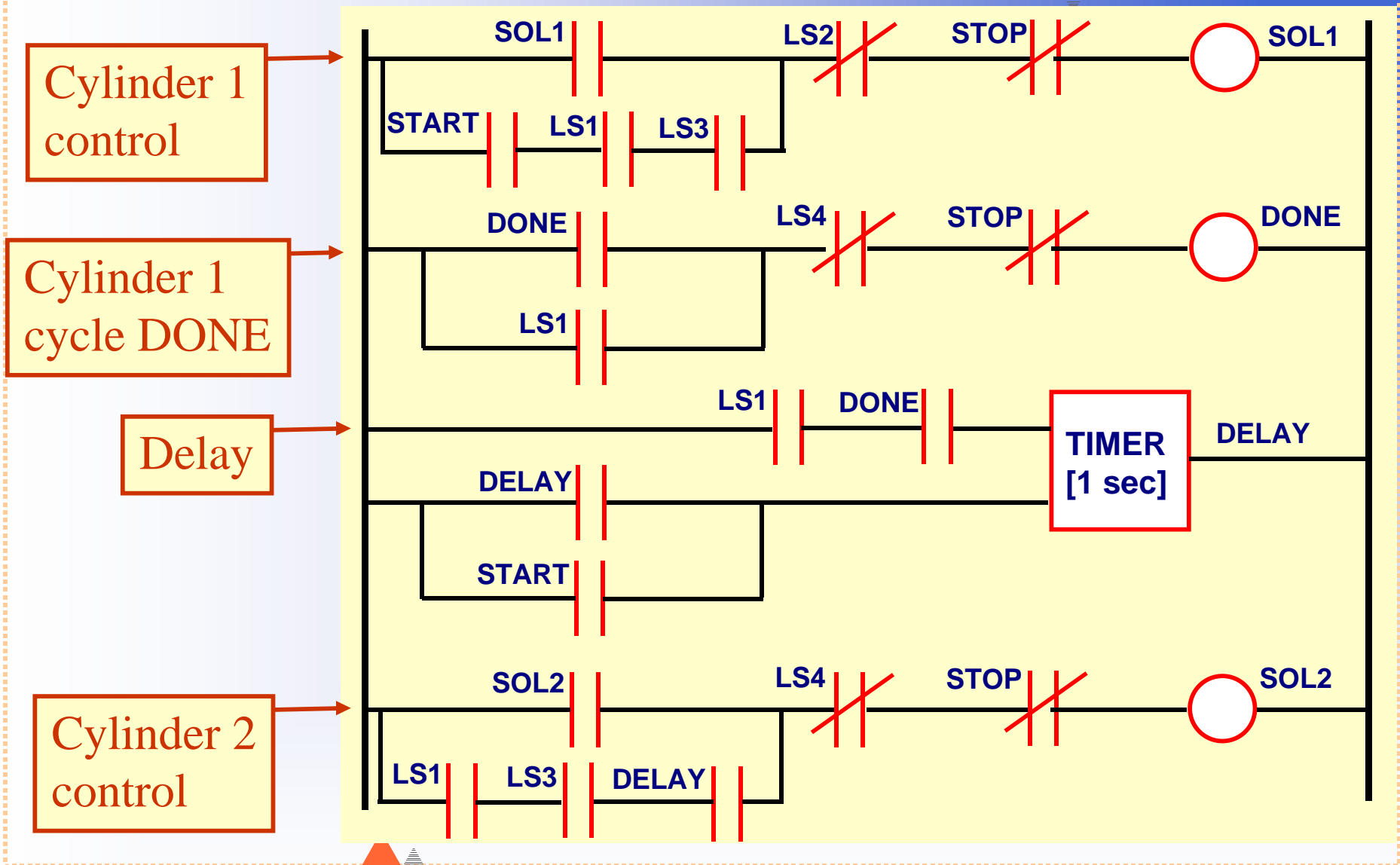
$$TIMER = LS1 \cdot DONE$$

CONDITION TO ACTIVATE
TIMER

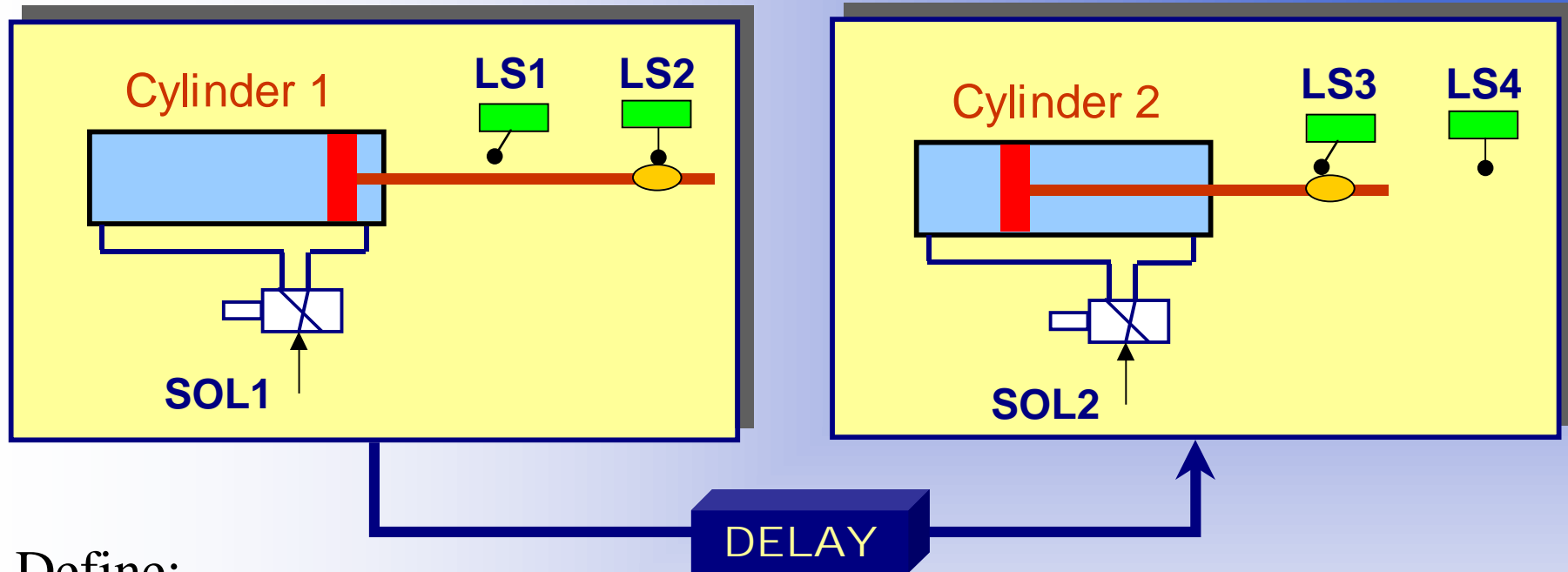
$$SOL2 = (SOL2 + LS1 \cdot LS3 \cdot \overline{DELAY}) \cdot \overline{LS4} \cdot \overline{STOP}$$



EXAMPLE - LADDER DIAGRAM



EXAMPLE – DES APPROACH



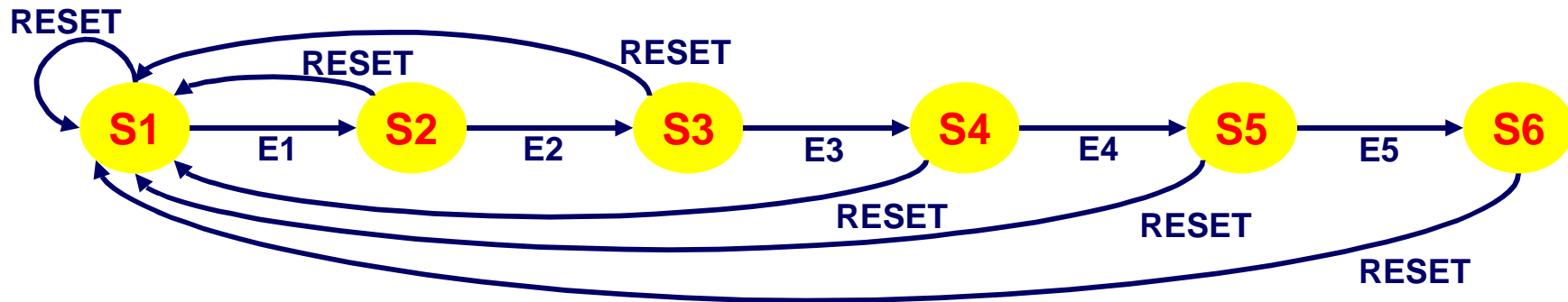
Define:

- System **STATES**
(e.g., “System at rest”, “Piston 1 moving right”)
- **EVENTS** causing state transitions
(e.g., “LS2 triggered”, “Timer activated”)



EXAMPLE - DES APPROACH

CONTINUED



STATES

| | |
|-----------|------------------------------------|
| S1 | Rest (LS1 , LS3 ON) |
| S2 | Piston 1 moves right |
| S3 | Piston 1 moves left |
| S4 | Timer ON |
| S5 | Piston 2 moves right |
| S6 | Piston 2 moves left |

EVENTS

| | |
|--------------|------------------------|
| E1 | START |
| E2 | LS2 triggered |
| E3 | LS1 triggered |
| E4 | TIMER activated |
| E5 | LS4 triggered |
| E6 | LS3 triggered |
| RESET | STOP |

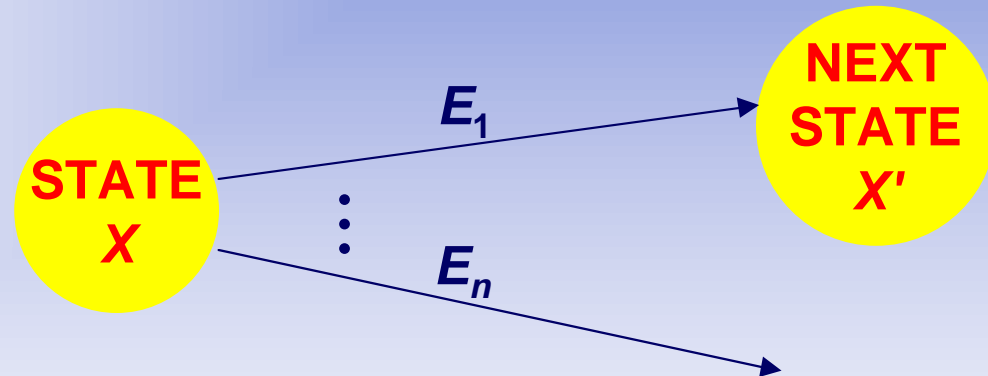


STATE TRANSITION LOGIC

Two main elements in a *Discrete Event System*:

1. **STATE** - the status of some system component
2. **EVENT** - instantaneous action causing a state transition

STATE TRANSITION
DIAGRAM:



STATE TRANSITION FUNCTION: $X' = f(X, E)$

NOTE: Only *FEASIBLE* events at state X are considered

DES MODELING: STATE AUTOMATON

AUTOMATON: (E, X, Γ, f, x_0)

E : Event Set

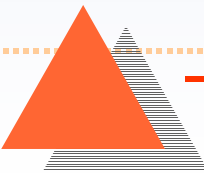
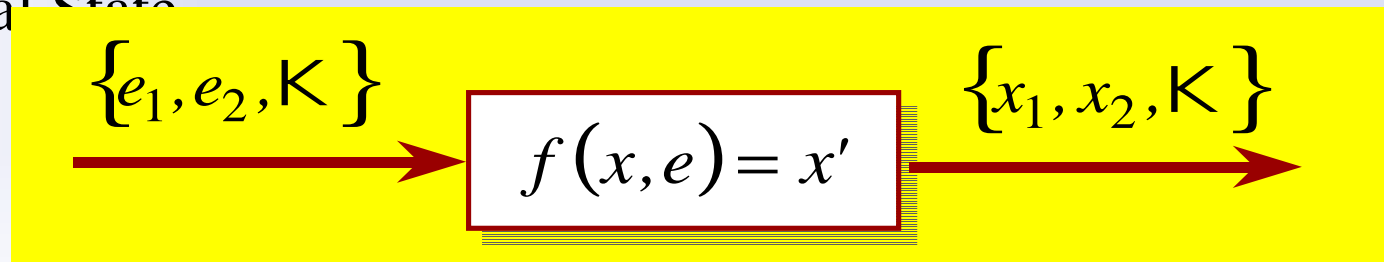
X : State Space

$\Gamma(x)$: Set of *feasible* or *enabled* events
at state x

$$f : X \times E \rightarrow X$$

f : State Transition Function
(undefined for $x_0 \in X$ and $e \notin \Gamma(x)$)

x_0 : Initial State

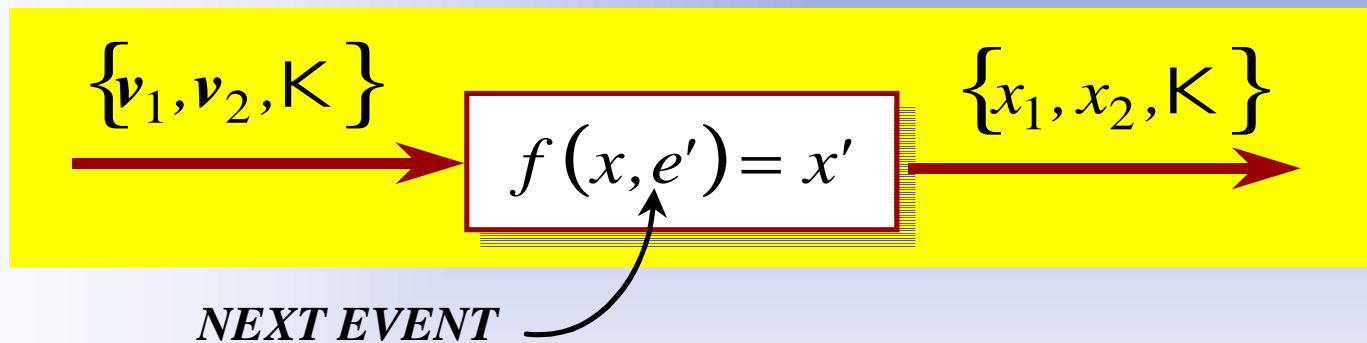


TIMED AUTOMATON

Add a *Clock Structure* V to the Automaton: $(E, X, \Gamma, f, x_0, V)$

where: $V = \{v_i : i \in E\}$

and v_i is a *Clock or Lifetime* Sequence: $v_i = \{v_{i1}, v_{i2}, \dots, \infty\}$
one for each event i



Need an *internal mechanism* to determine
NEXT EVENT e' and hence
NEXT STATE $x' = f(x, e')$



HOW THE TIMED AUTOMATON WORKS...

➔ CURRENT STATE

$x \in X$ with feasible event set: $\Gamma(x)$

➔ CURRENT EVENT

e that caused transition into x

➔ CURRENT EVENT TIME

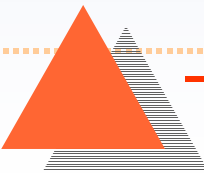
t associated with e



Associate a

CLOCK VALUE/RESIDUAL LIFETIME y_i

with each feasible event $i \in \Gamma(x)$



HOW THE TIMED AUTOMATON WORKS...

CONTINUED

➔ **NEXT/TRIGGERING EVENT e' :**

$$e' = \arg \min_{i \in \Gamma(x)} \{y_i\}$$

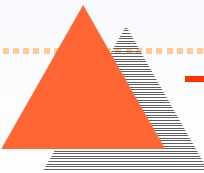
➔ **NEXT EVENT TIME t' :**

$$t' = t + y^*$$

$$\text{where: } y^* = \min_{i \in \Gamma(x)} \{y_i\}$$

➔ **NEXT STATE x' :**

$$x' = f(x, e')$$



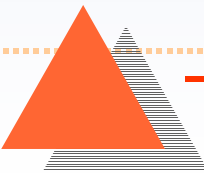
HOW THE TIMED AUTOMATON WORKS...

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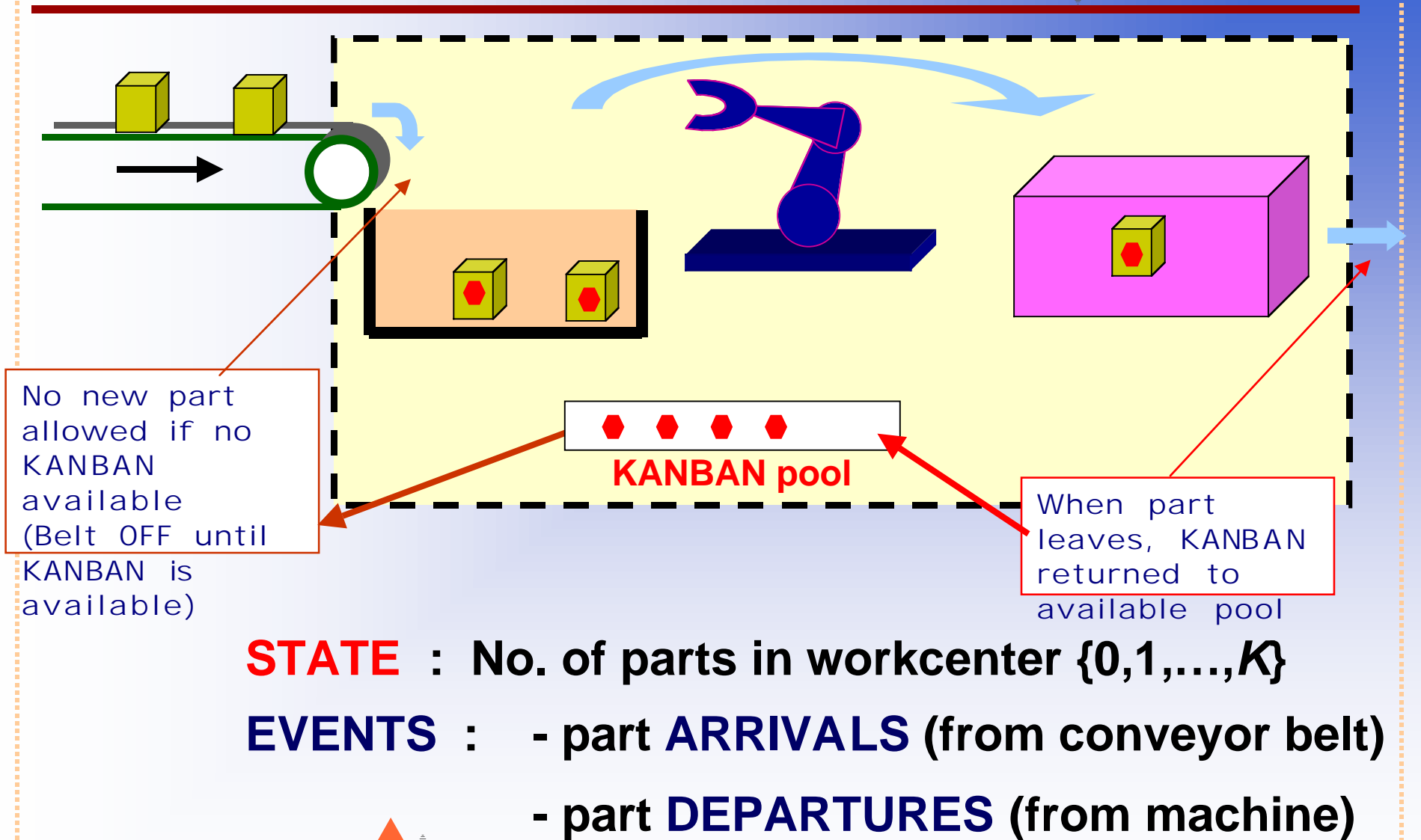
➔ Determine new **CLOCK VALUES** y'_i
for every event $i \in \Gamma(x)$

$$y'_i = \begin{cases} y_i - y^* & i \in \Gamma(x'), i \in \Gamma(x), i \neq e' \\ v_{ij} & i \in \Gamma(x') - \{\Gamma(x) - e'\} \\ 0 & \text{otherwise} \end{cases}$$

where: v_{ij} = new lifetime for event i

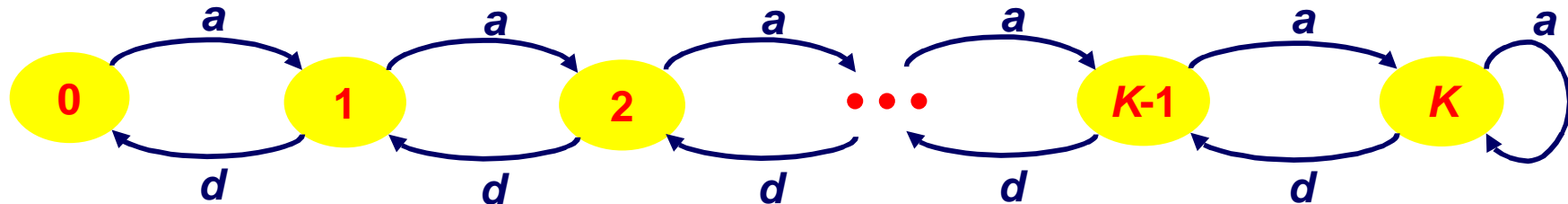


TIMED AUTOMATON - AN EXAMPLE



TIMED AUTOMATON - AN EXAMPLE

CONTINUED



$$E = \{a, d\} \quad X = \{0, 1, \dots, K\}$$
$$\Gamma(x) = \{a, d\} \quad \text{for all } x > 0$$
$$\Gamma(0) = \{a\}$$

$$f(x, e') = \begin{cases} x+1 & e' = a, x < K \\ x-1 & e' = d, x > 0 \end{cases}$$

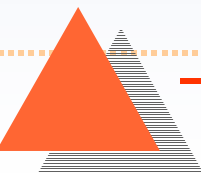
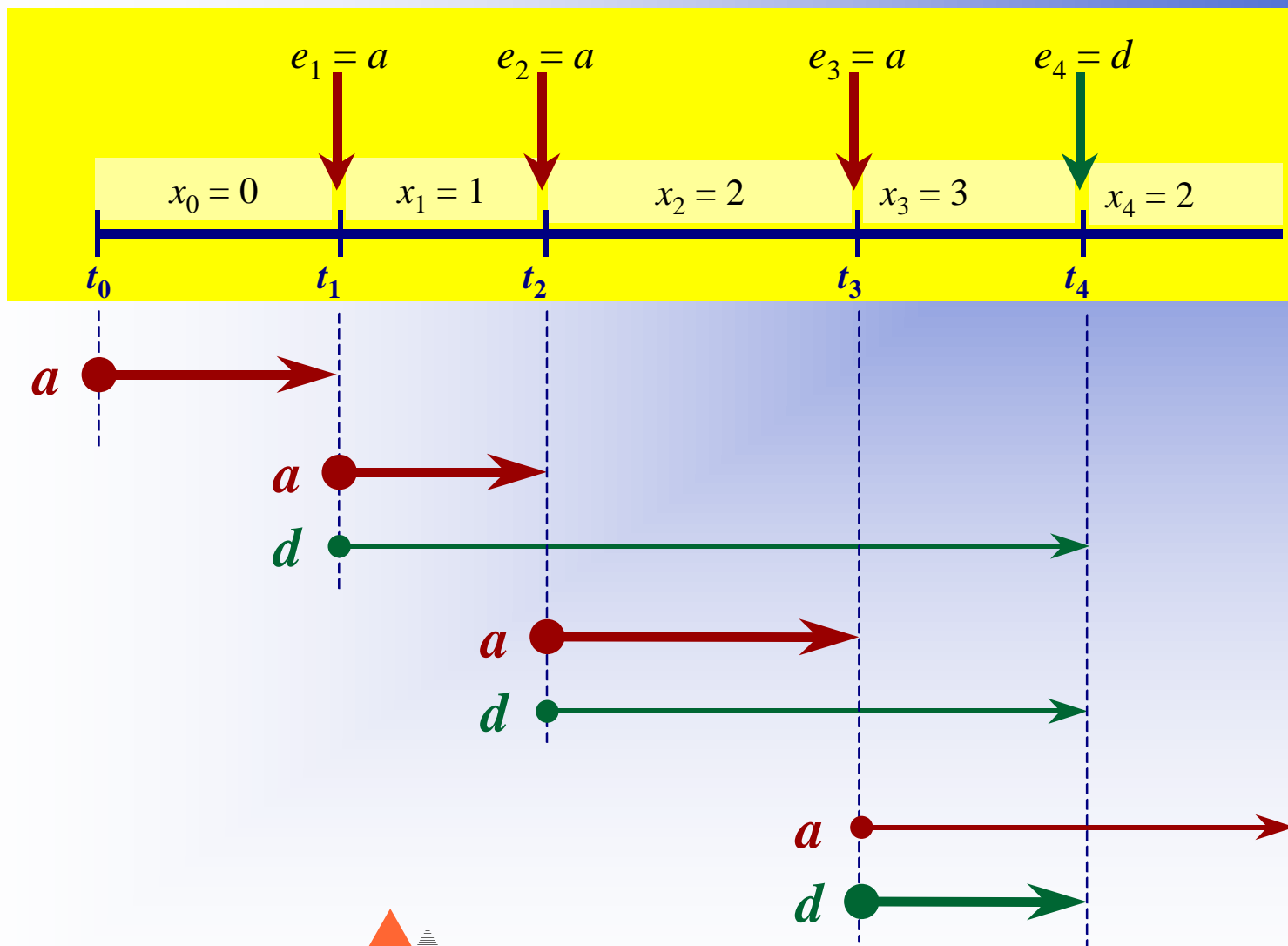
$$\text{Given input: } \nu_a = \{\nu_{a1}, \nu_{a2}, \dots, \nu_{aK}\} \quad \nu_d = \{\nu_{d1}, \nu_{d2}, \dots, \nu_{dK}\}$$



TIMED AUTOMATON - A TYPICAL SAMPLE PATH



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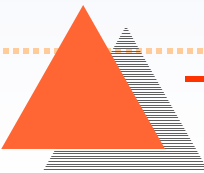
STOCHASTIC TIMED AUTOMATON

- Same idea with the Clock Structure consisting of *Stochastic Processes*
- Associate with each event i a *Lifetime Distribution* based on which ν_i is generated

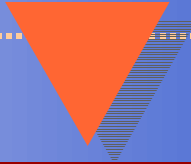


Generalized Semi-Markov Process
(GSMP)

- In a simulator, ν_i is generated through a pseudorandom number generator



UNTIMED DES: SUPERVISORY CONTROL



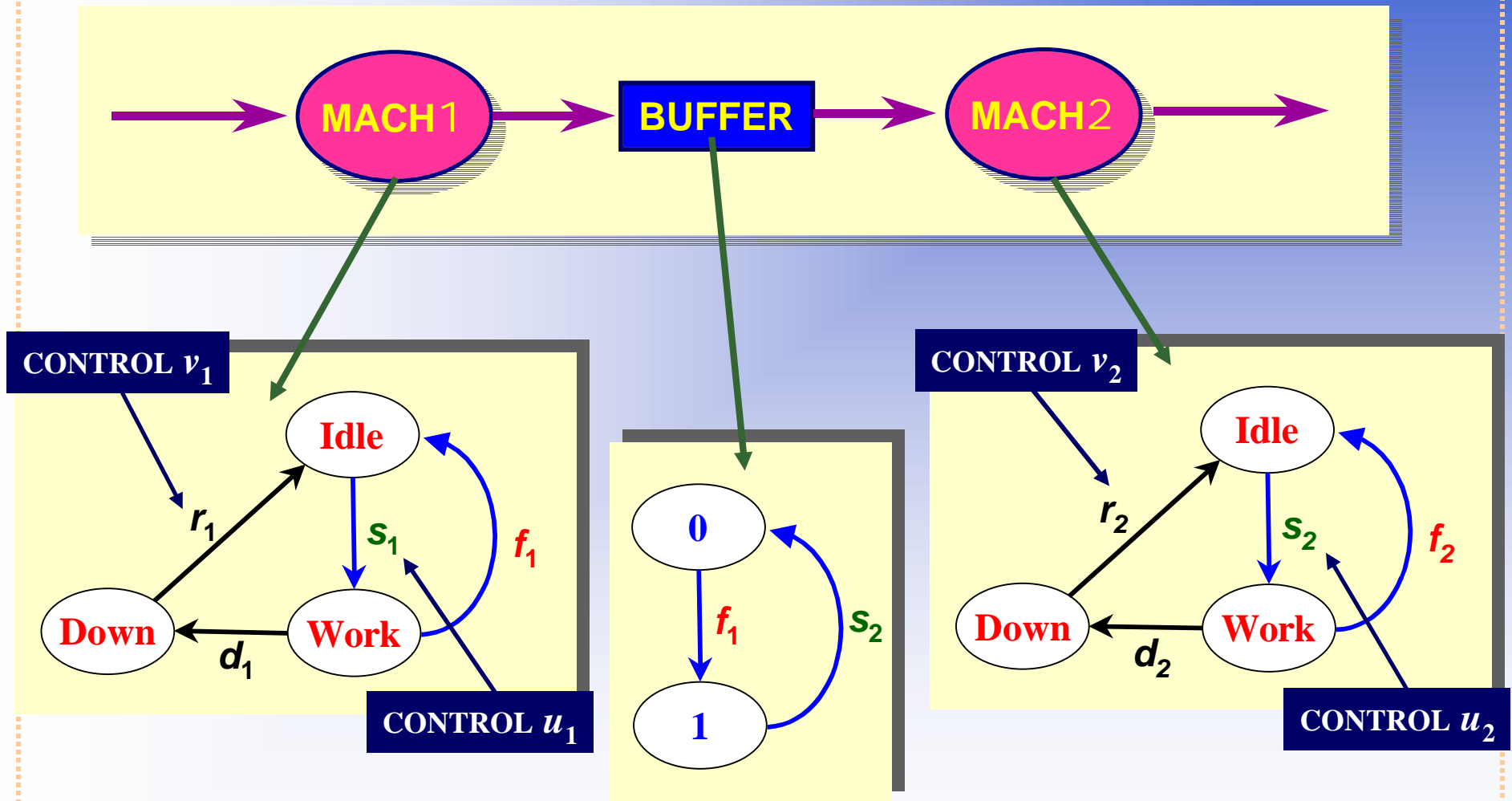
- Events may be CONTROLLABLE or UNCONTROLLABLE
- *Supervisory control:*

ENABLE/DISABLE controllable events so as to meet desired specifications (avoid deadlock, illegal states, etc.)

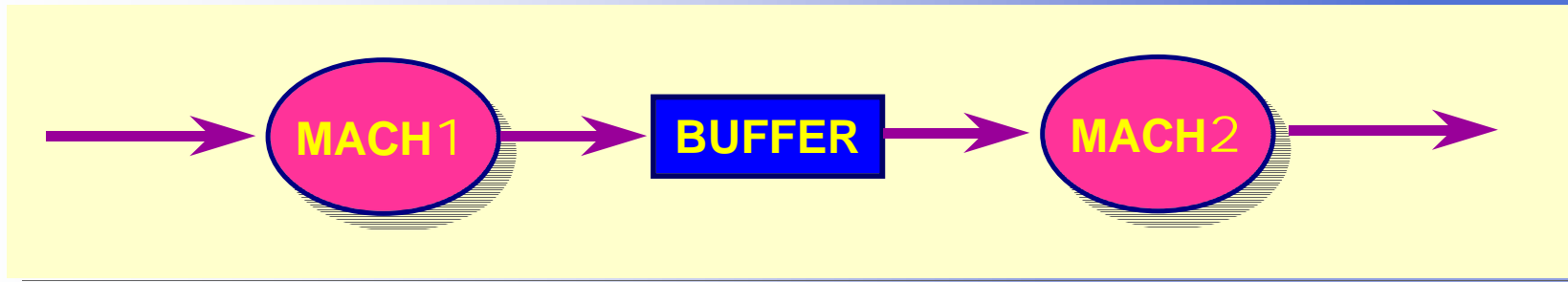
In the simplest case, all events are assumed *observable*.



SUPERVISORY CONTROL -- EXAMPLE



OPERATING RULES (SPECIFICATIONS)

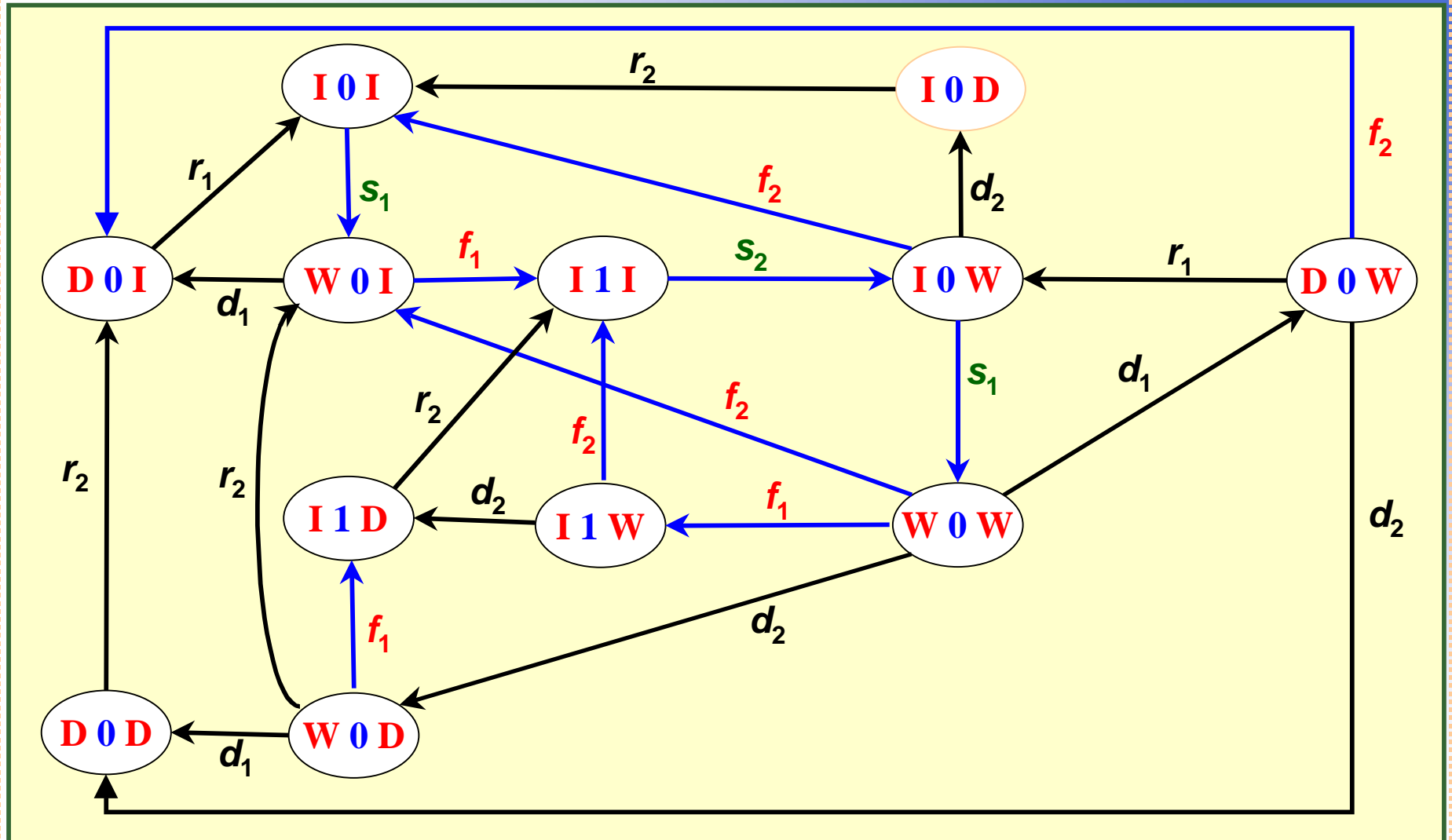


- MACH1 can only start when BUFFER is **empty**.
- MACH2 can only start when BUFFER is **full**.
- MACH1 cannot start when MACH2 is **down**.
- If both MACH1 and MACH2 are **down**, then MACH2 is repaired first



STATE TRANSITION DIAGRAM

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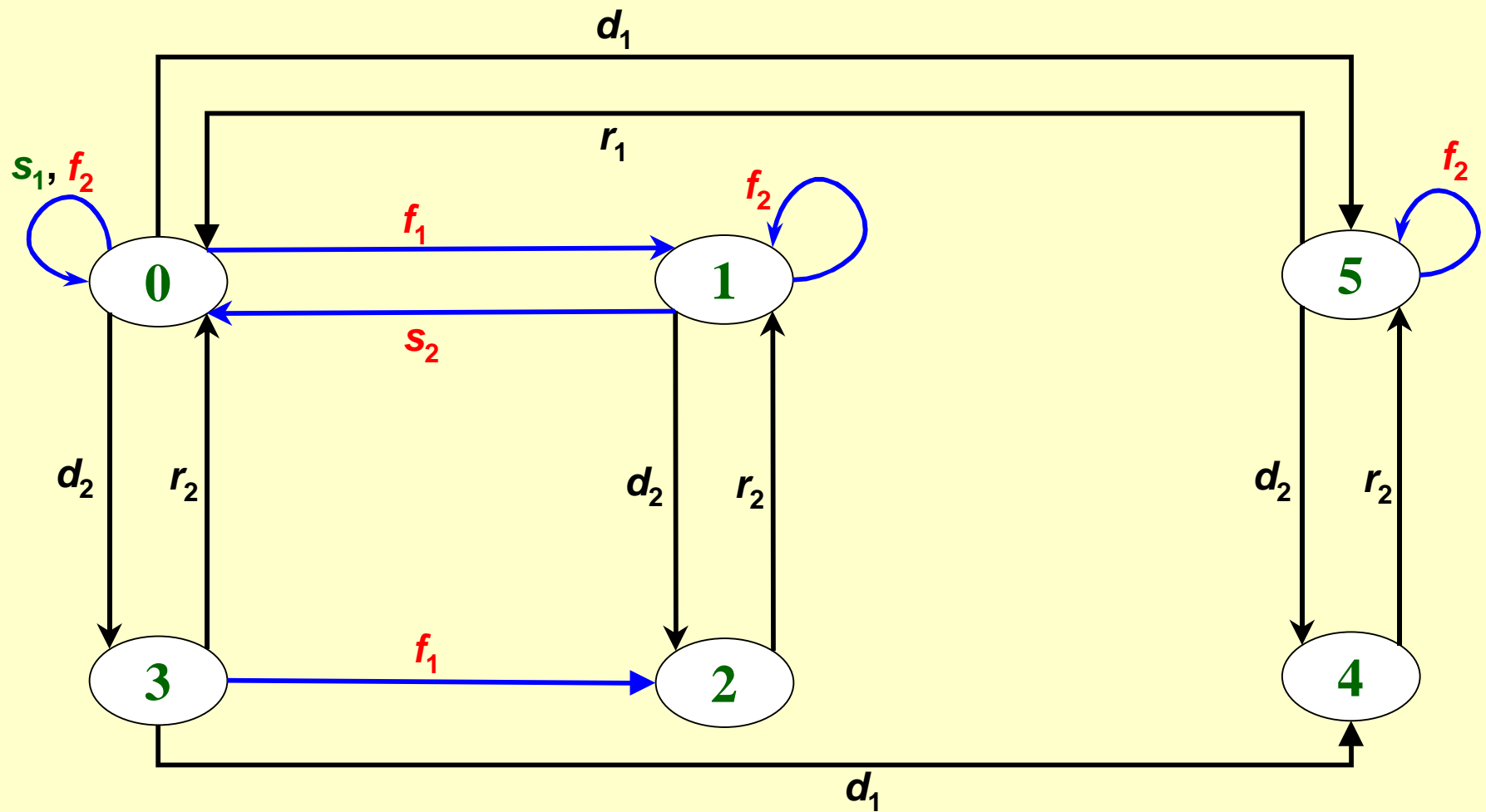
SUPERVISOR DESIGN

| STATE CONTROL | I0I | W0I | I1I | I0W | I0D | W0W | I1W | I1D | W0D | D0D | D0I | D0W |
|---------------|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|
| u_1 | 1 | 1 | 0 | 1 | 0 | 1 | 0 | 0 | 0 | 0 | - | - |
| v_1 | - | - | - | - | - | - | - | - | - | 0 | 1 | 1 |
| u_2 | 0 | 0 | 1 | 0 | 0 | 0 | 1 | - | 0 | 0 | 0 | 0 |
| v_2 | - | - | - | - | 1 | - | - | 1 | 1 | 1 | - | - |
| RED. STATE | 0 | 0 | 1 | 0 | 3 | 0 | 1 | 2 | 3 | 4 | 5 | 5 |

- 1 = ENABLE event
- 0 = DISABLE event
- = DON'T CARE

From a SUPERVISOR perspective, model is reduced to 6 states !

THE "QUOTIENT" SUPERVISOR



TIMED DES: SIMULATION AND 'RAPID LEARNING'



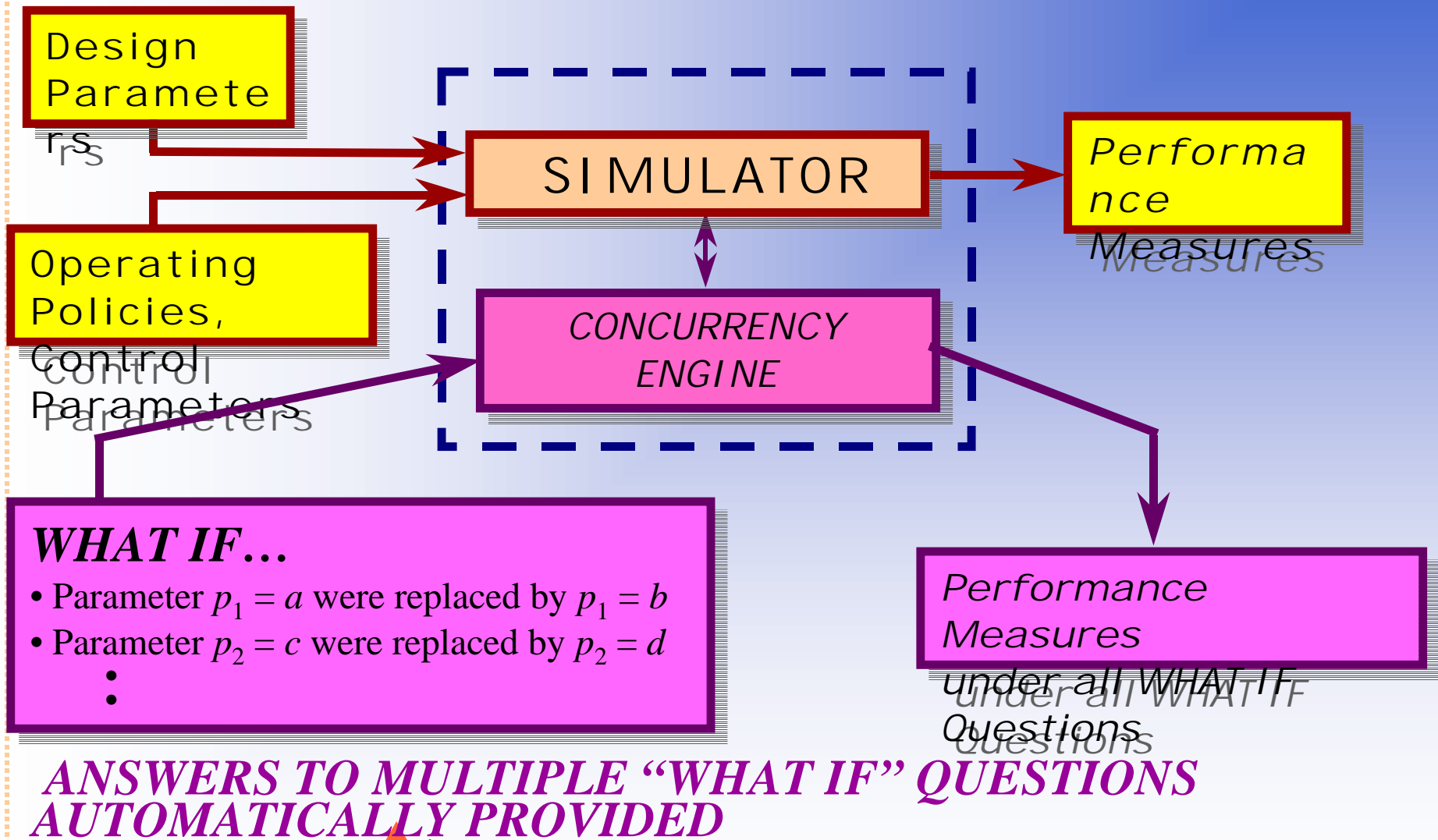
CONVENTIONAL SIMULATION ANALYSIS

- Repeatedly change parameters/operating policies
- Test different conditions
- Answer multiple WHAT IF questions

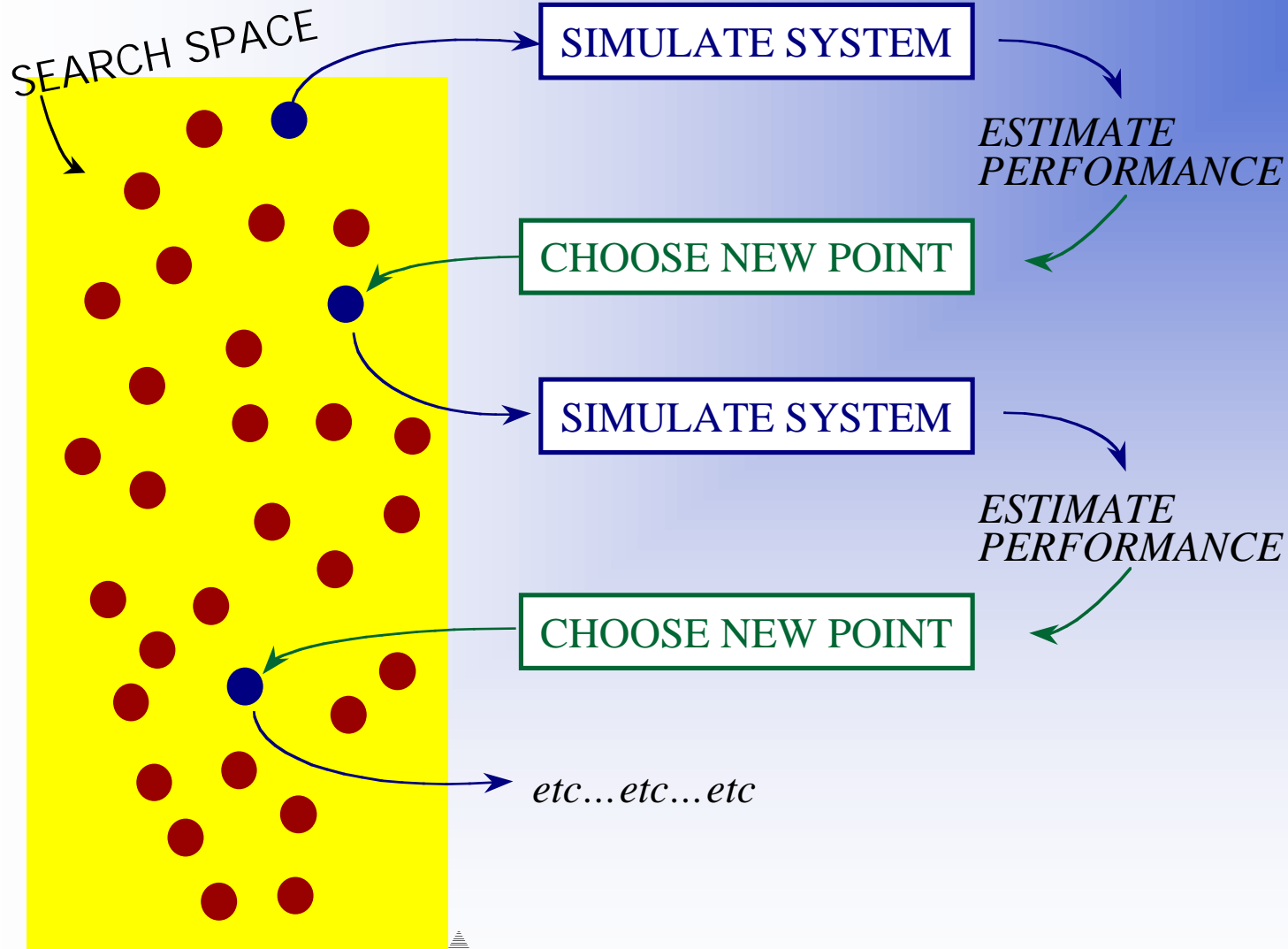


CONCURRENT SIMULATION

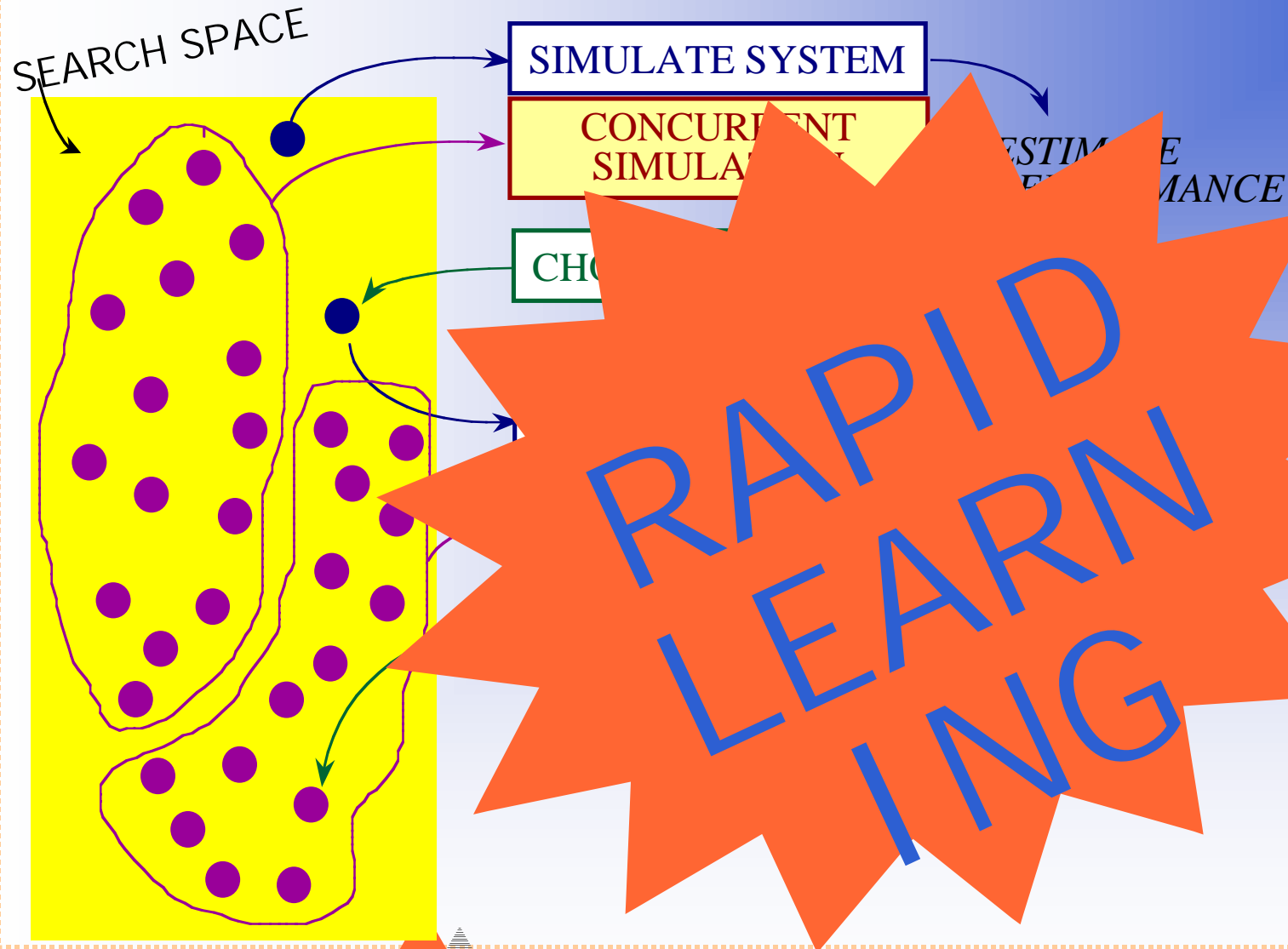
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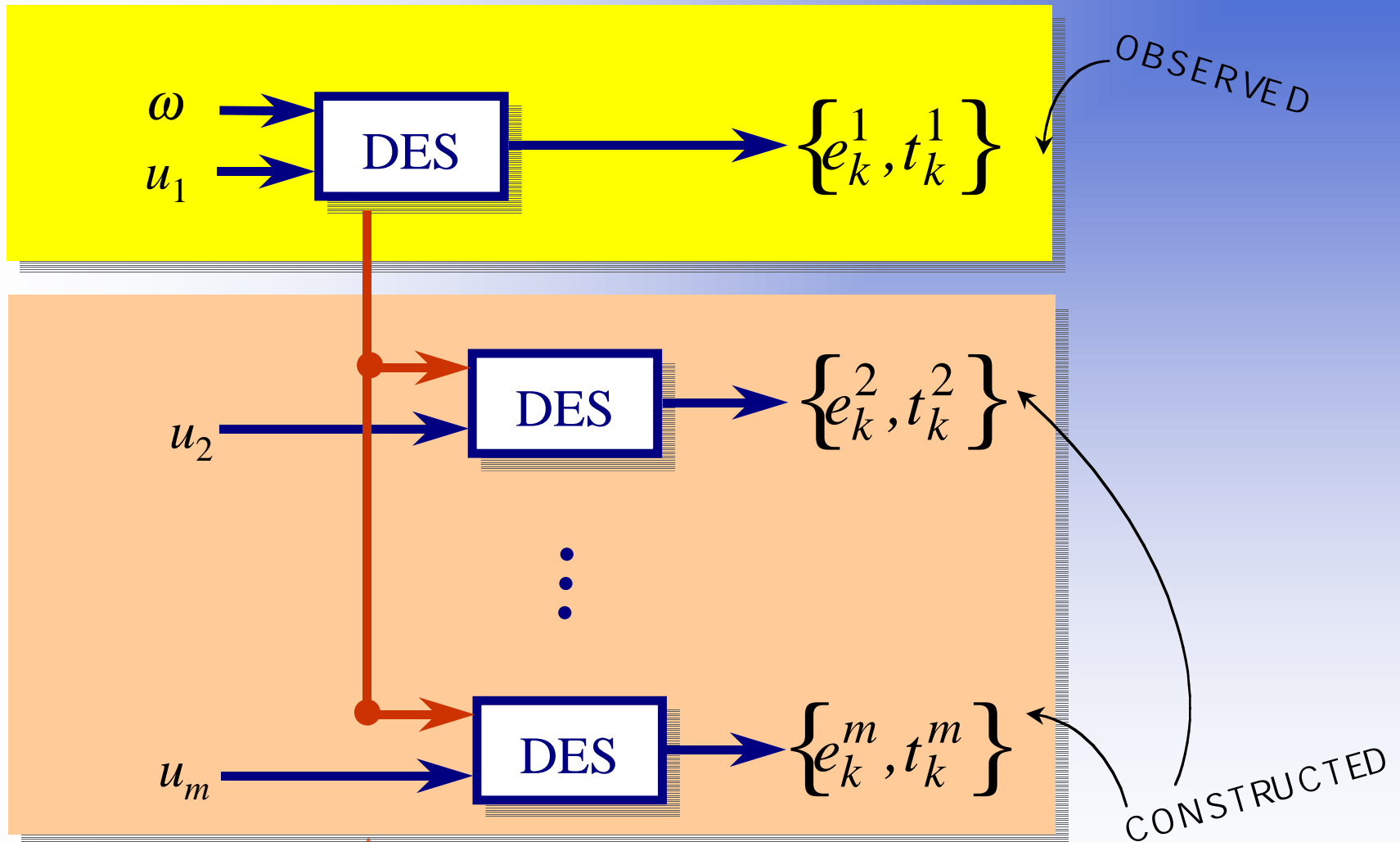
TYPICAL OPTIMIZATION OF COMPLEX SYSTEMS: *REPEATED TRIAL AND ERROR*



CONCURRENT SIMULATION FOR OPTIMIZATION



THE CONSTRUCTABILITY PROBLEM



CONSTRUCTABILITY: AN EXAMPLE



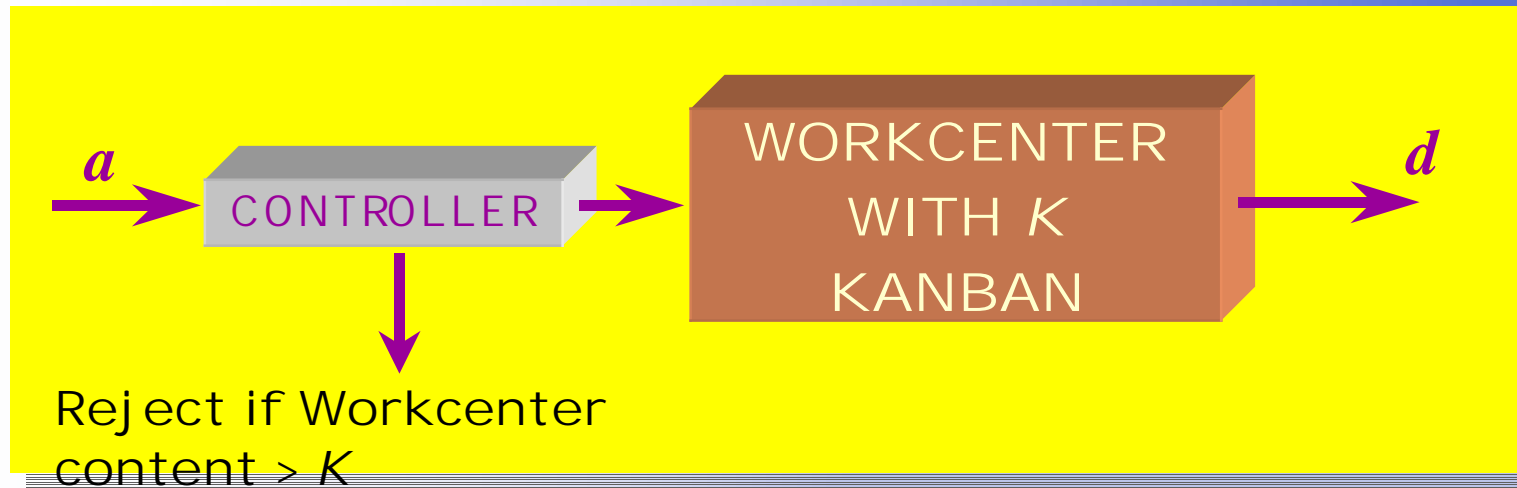
PROBLEM:

- How does the system behave under different choices of K ?
- What is the *optimal* K ?



CONSTRUCTABILITY: AN EXAMPLE

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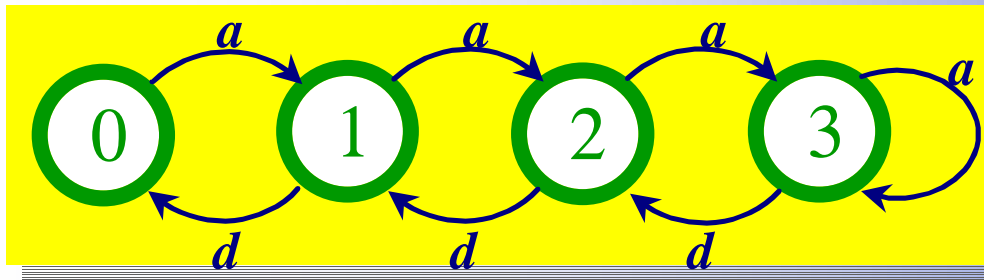
CONCURRENT SIMULATION APPROACH:

- Choose any K
- Simulate (or observe *actual system*) under K
- Apply Concurrent Simulation to **LEARN** effect of all other feasible K

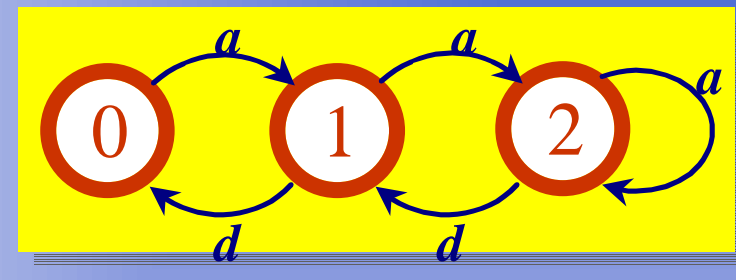


CONSTRUCTABILITY: AN EXAMPLE

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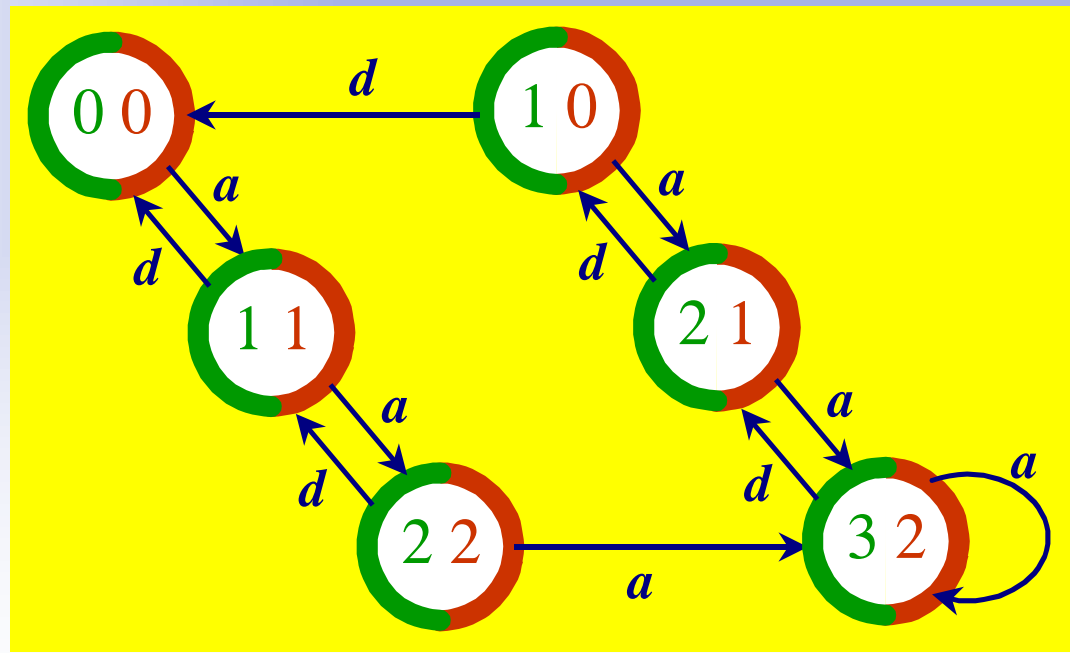


NOMINAL SYSTEM: $K = 3$



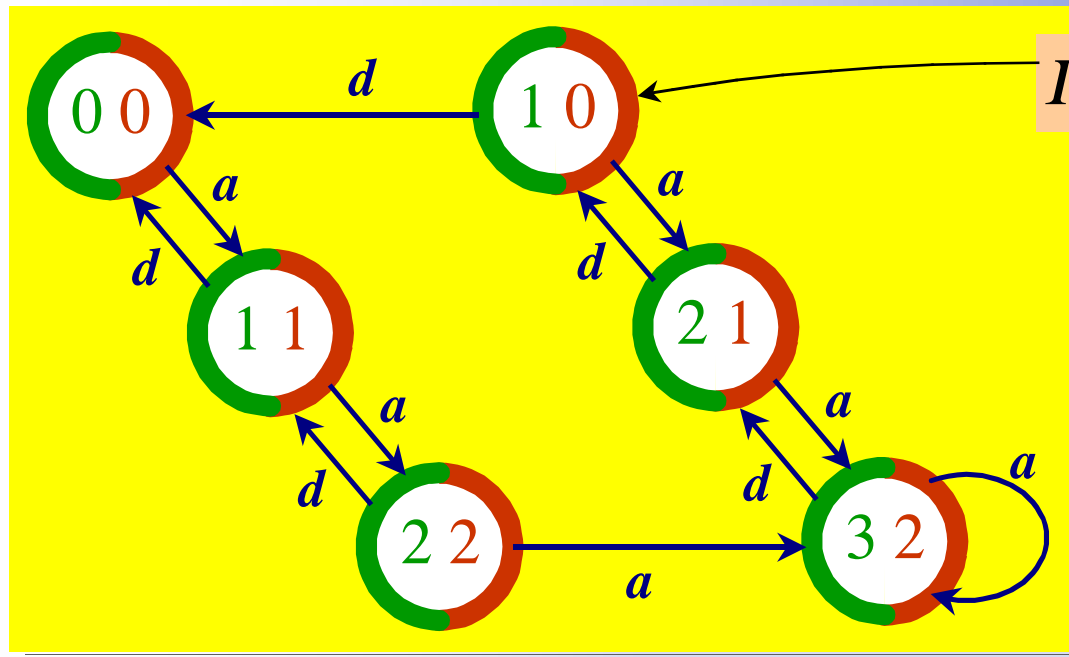
PERTURBED SYSTEM: $K = 2$

AUGMENTED SYSTEM



CONSTRUCTABILITY: AN EXAMPLE

CONTINUED



$$\Gamma(0) = \{a\} \subset \Gamma(1) = \{a, d\}$$

OBSERVABILITY
CONDITION

However, if roles of **NOMINAL** and **PERTURBED** are reversed, then things get a little trickier...

